#### CS442/542

- Minimize the DFA (D,  $\Sigma$ ,  $\delta$ , d<sub>0</sub>, D<sub>A</sub>}
- The algorithm builds a new machine from subsets of states of the original machine
- The algorithm first builds two subsets: the set of final states and the set of non-final states
- A subset is split if the subset has a conflict on a symbol
- A subset has a conflict on a symbol, c, when the transitions on c of two (or more) states in the subset do not go to states in the same subset.
- The algorithm halts when no subsets have conflicts (i.e. no more splits need to be done)

# Split (S is a set of states from the original DFA)

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\begin{array}{l} \mathsf{T} = \{ \ \mathsf{D}_{\mathsf{A}}, \ (\mathsf{D} - \mathsf{D}_{\mathsf{A}}) \ \};\\ \mathsf{P} = \varnothing\\ \text{while} \ ( \ \mathsf{P} \neq \mathsf{T} \ ) \ \mathsf{do}\\ & \mathsf{P} \leftarrow \mathsf{T};\\ & \mathsf{T} \leftarrow \varnothing;\\ & \mathsf{for} \ \mathsf{each} \ \mathsf{p} \in \mathsf{P} \ \mathsf{do}\\ & \mathsf{T} = \mathsf{T} \cup \mathsf{Split}(\mathsf{p});\\ & \mathsf{end};\\ \mathsf{end}; \end{array}
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- To build a minimized DFA from a DFA create a new DFA consisting of maximal consistent states which are subsets of states of the original DFA. These subsets form the set of states in the minimized DFA
- Consistent State
  - Suppose S is a set of states s1, s2,...Sn
  - Suppose  $\delta$  is the transition function of the original DFA
  - A set of states, S, is consistent if for each a in  $\Sigma$  and for each si in S,  $\delta(si,a) = qi$  where all the qi belong to the same consistent state
- A maximal consistent state is a state that is consistent and to which a new state cannot be added and the state remains consistent

- To produce a set of maximal consistent states do the follow
  - Create two subsets of states: final and non-final states
  - For each a in  $\Sigma$  if for some S all the transitions on a do not lead to the same subset of states, create subsets of S that lead to the same subset of states.
  - Continue the process until for each S in the new machine and for each a in  $\Sigma$  the transitions on a lead to the same subset

#### DFA for (0\*10\*10\*)\*



 Round 1 S0, S4, S5 S1, S2, S3 S1, S5, S5 **S1**, **S3**, **S3** 0 • Round 2 S0 S2 S4, S5 S1, S2, S3 S2, S2 S2, S4, S4 • Round 3 S4, S5 **S1** S2, S3 S0 S1 S2 S1 S2 S5, S5 S2, S2 S3, S3 S4, S4 0





### DFA for 0\*(10\*10\*)\*



- Round 1
  - S0, S1, S4, S50S1, S1, S5, S51S2, S2, S2, S2

S2, S3 S3, S3 S4, S4



- Potential Problems
  - The transition function is not total
  - Dead states