## CS 442/542

Parsing Part 3

## First Set

- For any grammar symbol $\alpha$, First( $\alpha$ ) is the set of terminals or epsilon that can appear at the start of a sentence derived from a


## Follow Set

- For any nonterminal $a$, Follow(a) is the set of terminals that can occur immediately after $\alpha$ in a sentence


## First ${ }^{+}$

- First ${ }^{+}(A->B)=\{t \mid$ if $\varepsilon \notin \operatorname{First}(B)$ then $t \in \operatorname{First}(B)$ else $t \in \operatorname{First}(B) \mathrm{U}$ Follow(A) \}
- Also called the Predict set
- Note this definition includes $\varepsilon$ in the First ${ }^{+}$set. Some definitions exclude it and in the way this set will be used in the algorithm to build the $\mathrm{LL}(1)$ parse table $\varepsilon$ will not be important


## Build First Set

```
for each }\alpha\in(T\cupeof\cup\epsilon) do
    FIRST(\alpha)}\leftarrow\alpha
end;
for each A \in NT do;
    FIRST(A) \leftarrow\emptyset;
end;
while (FIRST sets are still changing) do;
    for each p\inP, where p has the form }A->\beta\mathrm{ do;
        if \beta}\mathrm{ is }\mp@subsup{\beta}{1}{}\mp@subsup{\beta}{2}{}\ldots\mp@subsup{\beta}{k}{}\mathrm{ , where }\mp@subsup{\beta}{i}{}\inT\cupNT\mathrm{ , then begi,
                rhs }\leftarrow\operatorname{FIRST}(\mp@subsup{\beta}{1}{})-{\epsilon}
                i}\leftarrow1
                while ( }\epsilon\in\operatorname{FIRST}(\mp@subsup{\beta}{i}{})\mathrm{ and }i\leqk-1) do
                    rhs \leftarrowrhs U(FIRST ( }\mp@subsup{\beta}{i+1}{})-{\epsilon})
                    i}\leftarrowi+1
                end;
            end;
            if i = k and }\epsilon\in\operatorname{FIRST}(\mp@subsup{\beta}{k}{}
                then rhs \leftarrow rhs \cup {\epsilon};
            FIRST(A) \leftarrow FIRST(A) U rhs;
    end;
end:
```


## Build the First Sets for the Following Grammar

1. Goal -> E
2. E-> T E1
3. E1 -> + TE1
4. E1 -> - T E1
5. E1 $\rightarrow \varepsilon$
6. T -> F T1
7. T1 -> *FT1
8. T1 -> / F T1
9. $\mathrm{T} 1->\varepsilon$
10. $F$-> (E)
11. $F$-> num
12. F -> name

First(Goal) $=($, num, name
First(E) $=($, num, name
First(E1) $=+,-, \varepsilon$
First( $T$ ) $=($, num, name
First(T1) $=$ *, $/, \varepsilon$
First(F) $=($, num, name

## Build the First Sets for the Following Grammar

1. Goal -> E
2. E-> T E1
3. E1 -> + TE1
4. E1 -> - T E1
5. E 1 -> $\varepsilon$
6. T -> F T1
7. T1 -> *FT1
8. T1 -> / F T1
9. $\mathrm{T} 1->\varepsilon$
10. $F$-> (E)
11. $F$-> num
12. F -> name

First(Goal) $=$ First(E)
First( E ) $=$ First( T )
First(E1) $=+,-, \varepsilon$
First(T) $=$ First(F)
First(T1) $=$ *, $/, \varepsilon$
First(F) $=($, num, name

## Build the First Sets for the Following Grammar

1. Goal -> E
2. E-> T E1
3. E1 -> + TE1
4. E1 -> - T E1
5. E1 -> $\varepsilon$
6. T -> F T1
7. T1 -> *FT1
8. T1 -> / F T1
9. $\mathrm{T} 1->\varepsilon$
10. $F$-> (E)
11. $F$-> num
12. F -> name

First(Goal) $=$ First(E)
First( E ) $=$ First( T )
First(E1) $=+,-, \varepsilon$
First $(\mathbf{T})=$ First $(\mathrm{F})$
= (, num, name
First(T1) $=*, /, \varepsilon$
First(F) $=($, num, name

## Build the First Sets for the Following Grammar

1. Goal -> E
2. E-> T E1
3. E1 -> + TE1
4. E1 -> - T E1
5. E1 -> $\varepsilon$
6. T -> FT1
7. T1 -> *FT1
8. T1 -> / F T1
9. $\mathrm{T} 1->\varepsilon$
10. $F$-> (E)
11. $F$-> num
12. F -> name

First(Goal) $=$ First(E)
First $(\mathrm{E})=\operatorname{First}(\mathrm{T})$
= (, num, name
First(E1) $=+,-, \varepsilon$
First(T) $=\operatorname{First}(\mathrm{F})$
= (, num, name
First(T1) $=*, /, \varepsilon$
First(F) $=($, num, name

## Build the First Sets for the Following Grammar

1. Goal -> E
2. E-> T E1
3. E1 -> + TE1
4. E1 -> - T E1
5. E 1 -> $\varepsilon$
6. T -> F T1
7. T1 -> *FT1
8. T1 -> / F T1
9. T1 -> $\varepsilon$
10. $F$-> (E)
11. $F$-> num
12. F -> name

First(Goal) $=$ First(E)
= (, num, name
First(E) $=\operatorname{First}(T)$
= (, num, name
First(E1) $=+,-, \varepsilon$
First(T) = First(F)

$$
=(, \text { num, name }
$$

First(T1) $=$ *, $/, \varepsilon$
First(F) $=($, num, name

## Build Follow Sets

```
for each A\inNT do;
    FOLLOW(A) }\leftarrow\emptyset
end;
FOLLOW(S) \leftarrow {eof};
while (Follow sets are still changing) do:
        for each }p\inP\mathrm{ of the form }A->\mp@subsup{\beta}{1}{}\mp@subsup{\beta}{2}{}\cdots\mp@subsup{\beta}{k}{}\mathrm{ do;
        TRAILER }\leftarrow\operatorname{FOLLOW(A);
            for }1\leftarrowk\mathrm{ down to 1 do;
            if }\mp@subsup{\beta}{i}{}\inNT\mathrm{ then begin;
                FOLLOW ( }\mp@subsup{\beta}{i}{})\leftarrow\operatorname{FOLLOW}(\mp@subsup{\beta}{i}{})\cup\mathrm{ TRAILER;
                    if }\epsilon\in\operatorname{FIRST}(\mp@subsup{\beta}{i}{}
                        then TRALLER }\leftarrowT\mathrm{ Trailer U (FIRST( }\mp@subsup{\beta}{i}{})-\epsilon)
                        else Trailer }\leftarrow\operatorname{FIRST}(\mp@subsup{\beta}{i}{})\mathrm{ ;
                end;
                else TRAILER }\leftarrow\operatorname{FIRST}(\mp@subsup{\beta}{i}{});\quad// is {\mp@subsup{\beta}{i}{}
            end;
        end;
    end;
```


## Build the Follow Sets for the Following Grammar

1. Goal -> E
2. E-> T E1
3. E1 -> + T E1
4. E1 -> - T E1
5. E1 -> $\varepsilon$
6. T->FT1
7. T1 -> *FT1
8. T1 -> / FT1
9. T1 -> $\varepsilon$
10. $F$-> (E)
11. $F$-> num
12. F -> name

Follow(Goal) = eof<br>Follow(E) = eof, )<br>Follow(E1) = eof, )<br>Follow(T) = +, -, eof, )<br>Follow(T1) = +, -, eof, )<br>Follow(F) = *, /, +, - , eof, )

## Build the Follow Sets for the Following Grammar

1. Goal -> E
2. E-> T E1
3. E1 -> + T E1
4. E1 -> - T E1
5. E1 -> $\varepsilon$
6. T->FT1
7. T1 -> *FT1
8. T1 -> / FT1
9. T1 -> $\varepsilon$
10. $F$-> (E)
11. $F$-> num
12. F -> name

Follow (Goal) = eof<br>Follow(E) = Follow(Goal), )<br>Follow(E1) = Follow(E)<br>Follow(T) $=$ First(E1)<br>Follow(T1) = Follow(T)<br>Follow(F) = First(T1)

## Build the Follow Sets for the Following Grammar

1. Goal -> E
2. E-> T E1
3. E1 -> + T E1
4. E1 -> - T E1
5. E1 $\rightarrow \varepsilon$
6. T->FT1
7. T1 -> *FT1
8. T1 -> / FT1
9. T1 -> $\varepsilon$
10. $F$-> (E)
11. F -> num
12. F -> name


## Build the Follow Sets for the Following Grammar

1. Goal -> E
2. E-> T E1
3. E1 -> + T E1
4. E1 -> - T E1
5. E1 -> $\varepsilon$
6. T->FT1
7. T1 -> *FT1
8. T1 -> / FT1
9. T1 $->\varepsilon$
10. $F$-> (E)
11. F -> num
12. F -> name

$$
\begin{aligned}
& \text { Follow(Goal) = eof } \\
& \text { Follow(E) = Follow(Goal), ) } \\
& \text { = eof, ) } \\
& \text { Follow(E1) }=\text { Follow(E) } \\
& \text { = eof, ) } \\
& \text { Follow(T) }=\operatorname{First}(\mathrm{E} 1) \\
& =+,- \text {, } \operatorname{Follow}(E 1) \text {, } \\
& \text { Follow(E) } \\
& \text { = +, -, eof, ) } \\
& \text { Follow(T1) = Follow(T) } \\
& =+ \text {, -, Follow(E1), } \\
& \text { Follow(E) } \\
& \text { = +, -, eof, ) } \\
& \text { Follow(F) }=\text { First(T1) } \\
& ={ }^{*}, / \text {, Follow(T), } \\
& \text { Follow(T1) } \\
& =*, l,+,-, \text { eof, ) }
\end{aligned}
$$

## Build LL(1) Parse Table

build FIRST, FOLLOW, and FIRST ${ }^{+}$sets;
for each nonterminal A do;
for each terminal $w$ do;
Table[A,W] $\leftarrow$ error;
end;
for each production $p$ of the form $A \rightarrow \beta$ do:
for each terminal $W \in \operatorname{FIRST}^{+}(A \rightarrow \beta)$ do;
Table[A,w] $\leftarrow p$;
end;
if eof $\in \operatorname{FIRST}^{+}(A \rightarrow \beta)$
then Table[A, eof] $\leftarrow p$;
end;
end;

## Build the LL(1) Parse Table for the Following Grammar

| 1. Goal -> E |  | eof | + | - |  | / | ( |  | num | name |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 2. E -> T E1 | G |  |  |  |  |  | 1 |  | 1 | 1 |
| 3. E1 -> + T E1 | E |  |  |  |  |  | 2 |  | 2 | 2 |
| 4. E1 -> - T E1 | E1 | 5 | 3 | 4 |  |  |  | 5 |  |  |
| 5. E1 $\rightarrow$ ¢ | T |  |  |  |  |  | 6 |  | 6 | 6 |
| 6. T -> F T1 | T1 | 9 | 9 | 9 | 7 | 8 |  | 9 |  |  |
| 7. T1 -> * F T1 | F |  |  |  |  |  | 10 |  | 11 | 12 |
| 8. T1 -> / F T1 |  |  |  |  |  |  |  |  |  |  |
| 9. T1 -> $\varepsilon$ |  |  |  |  |  |  |  |  |  |  |
| 10. $F$-> (E) |  |  |  |  |  |  |  |  |  |  |
| 11. $F$-> num |  |  |  |  |  |  |  |  |  |  |
| 12. F -> name |  |  |  |  |  |  |  |  |  |  |

## Build the First+ Sets for the Following Grammar

| 1. Goal -> E | $\operatorname{First}(\mathrm{E})=($, num, name |
| :---: | :---: |
| 2. E -> T E1 | First( T$)=$ (, num, name |
| 3. E1 -> + T E1 | + |
| 4. E1 ->-T E1 | - |
| 5. E1 $\rightarrow$ ¢ | Follow(E1) = eof, ) |
| 6. T -> F T1 | First(F) = (, num, name |
| 7. T1 -> * F T1 | * |
| 8. T1 -> / F T1 | / |
| 9. $\mathrm{T} 1->\varepsilon$ | Follow(T1) = +, -, eof, ) |
| 10. F -> (E) | ( |
| 11. $F$-> num | num |
| 12. F -> name | name |

## Build the LL(1) Parse Table for the Following Grammar

| 1. Goal -> E |  | eof | + | - |  | / |  |  | num name |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 2. E -> T E1 | G |  |  |  |  |  | 1 |  | 1 | 1 |
| 3. E1 -> + T E1 | E |  |  |  |  |  | 2 |  | 2 | 2 |
| 4. E1 -> - T E1 | E1 | 5 | 3 | 4 |  |  |  | 5 |  |  |
| 5. E1 $\rightarrow$ ¢ | T |  |  |  |  |  | 6 |  | 6 | 6 |
| 6. T -> F T1 | T1 | 9 | 9 | 9 | 7 | 8 |  | 9 |  |  |
| 7. T1 -> * F T1 | F |  |  |  |  |  | 10 |  | 11 | 12 |
| 8. T1 -> / F T1 |  |  |  |  |  |  |  |  |  |  |
| 9. T1 -> $\varepsilon$ |  |  |  |  |  |  |  |  |  |  |
| 10. F -> (E) |  |  |  |  |  |  |  |  |  |  |
| 11. $F \rightarrow>$ num |  |  |  |  |  |  |  |  |  |  |
| 12. F -> name |  |  |  |  |  |  |  |  |  |  |

# Build the First, Follow and First+ Sets and the LL(1) Parse Table for the Following Grammar 

1. $S \rightarrow B$
2. $\mathrm{B}->\mathrm{B} 1 \mathrm{~B} 2$
3. B2 -> OR B1 B2
4. B2 -> $\varepsilon$
5. B1 $\rightarrow$ B3 B4
6. B4 -> AND B3 B4
7. B4-> $\varepsilon$
8. B3 -> VAR
