

## Sorting Algorithms

- There are dozens of sorting algorithms ${ }^{1}$
- Sorting algorithms can be evaluated in many ways
- run time
- memory usage
- general approach
- e.g., exchanging, sorting
- parallelizability
-i.e., can it be performed in parallel?


## Sorting

- Sorted data is crucial to human usability of data
$\cdot$ e.g., phone book, entering grades, dates
- Sorted data is also crucial to computational efficiency in accessing data
- i.e., how can a computer most efficiently find data?
- So, how do we sort data?


## Sorting Algorithms We'll Explore

- selection sort (typically iterative)
- insertion sort (typically iterative)
- merge sort (typically recursive)
- quicksort (typically recursive)


## Selection Sort

- Considered one of the classic sorting algorithms
- Very simple, but very inefficient (this tradeoff often occurs)
-Thumbnail sketch:
- scan through the array multiple times
- each time find the smallest "remaining" element
- move that element to correct position


## Selection Sort

- Array is divided into two parts: sorted (left part) and unsorted (right part) - initially, everything is unsorted
- Scan through the unsorted part for the smallest element
- Swap the smallest element with the leftmost unsorted value
- Length of sorted part increases by one, length of unsorted part decreases by one
- Repeat


## Selection Sort

Array is divided into two parts: sorted (left part) and unsorted (right part)
Scan through the unsorted part for the smallest element
Swap the smallest element with the leftmost unsorted value

Length of sorted part increases by one, length of unsorted part
decreases by one
$\begin{array}{llllll}8 & 3 & 2 & 5 & 9 & 7\end{array}$

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$8 \quad 3 \quad 2 \quad 5 \quad 9 \quad 7$
smallestIndex = 1

## Selection Sort

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(left part) and unsorted (right part)
Scan through the unsorted part for the smallest element
Swap the smallest element with the
leftmost unsorted value
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decreases by one
$83 \quad 25 \quad 9 \quad 7$
smallestIndex = 2

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83259
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| Selection Sort |  |  |  |
| :--- | :--- | :--- | :--- | :--- |
| Array is divided into two parts: sorted <br> (left part) and unsorted (right part) <br> Scan through the unsorted part for <br> the smallest element <br> Swap the smallest element with the <br> leftmost unsorted value | 8 | 2 |  |
| Length of sorted part increases by <br> one, length of unsorted part <br> decreases by one |  |  |  |

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\end{array}
$$

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\section*{| 2 | 3 | 8 | 5 | 9 | 7 |
| :--- | :--- | :--- | :--- | :--- | :--- |}

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Array is divided into two parts: sorted (left part) and unsorted (right part)

Scan through the unsorted part for the smallest element

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Length of sorted part increases by

smallestIndex = 3

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| 2 | 3 | 8 | 5 | 9 | 7 |
| :--- | :--- | :--- | :--- | :--- | :--- |

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Scan through the unsorted part for the smallest element
Swap the smallest element with the leftmost unsorted value

Length of sorted part increases by one, length of unsorted part
decreases by one

```
2
smallestIndex = 5
```

| Selection Sort <br> Array is divided into two parts: sorted <br> (left part) and unsorted (right part) <br> Scan through the unsorted part for <br> the smallest element <br> Swap the smallest element with the <br> leftmost unsorted value <br> Length of sorted part increases by <br> one, length of unsorted part <br> decreases by one$\quad$2 3 5 9 |
| :--- |

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## Insertion Sort

- Considered one of the classic sorting algorithms
- Very simple, but very inefficient
-Thumbnail sketch:
- places next unsorted element into sorted part of array by...
-...searching for correct position within the sorted part
-that position may not be the element's final position


## Insertion Sort

- Array is divided into two parts: sorted (left part) and unsorted (right part) - initially, first element is sorted, everything else is unsorted
- Look at the leftmost unsorted value
- Move it down the sorted list until it is in the correct place
- Length of sorted part increases by one, length of unsorted part decreases by one
- Repeat


## Insertion Sort

Array is divided into two parts: sorted
(left part) and unsorted (right part)
Look at the leftmost unsorted value
Move it down the sorted list until it is in the correct place
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| 3 | 8 | 2 | 5 | 9 | 7 |
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## Insertion Sort

Array is divided into two parts: sorted
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Move it down the sorted list until it is in the correct place

```
3
```

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$$
\begin{array}{|lll|lll}
\hline 2 & 3 & 8 & 5 & 7
\end{array}
$$

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$$
\begin{array}{llllll}
2 & 3 & 5 & 7 & 8 & 9
\end{array}
$$

Move it down the sorted list until it is in the correct place

Length of sorted part increases by
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decreases by one

| 2 | 3 | 5 | 7 | 8 | 9 |
| :--- | :--- | :--- | :--- | :--- | :--- |



## Insertion Sort

Array is divided into two parts: sorted
(left part) and unsorted (right part)
Look at the leftmost unsorted value
$\qquad$

Length of sorted part increases by one, length of unsorted part
decreases by one

## Insertion Sort: Code

```
private static void insertionSort(int ar[]) {
    // Gradually look at each unsorted number
    for (int i = 0; i < arr.length; i++) {
        // Current value to sort
        int value = arr[i];
        int value = ,
        // Shift elements until value is in place
        while ( }j>=0\mathrm{ && arr[j] > value) {
            arr[j+1] = arr[j];
        } j-
    arr[j+1] = value;
}
```


## Selection Sort \& Insertion Sort

private static void selectionsort(int arr []) \{
// Gradually move boundary of unsorted
Gradually move boundary of unsorted portion
(nti = 0; i < arr. length-1; i++) \{
// Find the index of the smallest unsorted item
int smallestIndex $=\mathrm{i}$;
int smallestIndex $=\mathrm{i}$;
for (int $j=i+1 ; ~$
< arr.length; $j++$ ) $\{$
int $j=i+1 ; j$ < arr. length; $j++)$ \{
$f($ arr $[j]=$ arr[smallestIndex] $)\{$
smallestIndex $=j$;
\} $\}$
// Swap
int temp $=\operatorname{arr}[$ smallestIndex]; .
$\operatorname{arr}[$ smallestIndex] $=\operatorname{arr}[i]$;
$\operatorname{arr}[i]=\operatorname{temp}$;


## Algorithm Analysis

First, consider what is the best and worst case scenarios for sorting an array Then, fill out the chart below with the run time (i.e., big O ):

|  | selection sort | insertion sort |
| :---: | :---: | :---: |
| best case scenario |  |  |
| worst case scenario |  |  |

## Algorithm Analysis

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## Algorithm Analysis

First, consider what are the best and worst case scenarios for each algorithm Then, fill out the chart below with the run time (i.e., big O ):

|  | selection sort | insertion sort |
| :---: | :---: | :---: |
| best case scenario | $O\left(n^{2}\right)$ | $O(n)$ <br> (already sorted) |
| worst case scenario | $O\left(n^{2}\right)$ | $O\left(n^{2}\right)$ <br> (reverse order) |

## Merge Sort

Array is divided into its smallest unit
i.e., a single element

Sort and merge each paired subarray of elements
Repeat sort/merge until there is only one array

## Merge Sort

- Considered one of the classic sorting algorithms
- More complex than selection/insertion sort
-...but more efficient!
-Thumbnail sketch:
- break the array up into individual elements
- sorts pairs of elements, then pairs of pairs, etc...until you have one unified array


## Merge Sort

Array is divided into its smallest unit
$\begin{array}{llllllll}6 & 4 & 8 & 3 & 2 & 5 & 9 & 7\end{array}$ i.e., a single element

Sort and merge each paired subarray of elements
Repeat sort/merge until there is only one array

## Merge Sort

Array is divided into its smallest unit
i.e., a single element

```
|6
```

Sort and merge each paired subarray
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Repeat sort/merge until there is only
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## Merge Sort

Array is divided into its smallest unit
i.e., a single element

```
46 4 8 3 3 2 5 5 9 7
```

Sort and merge each paired subarray of elements

Repeat sort/merge until there is only
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## Merge Sort

```
Array is divided into its smallest unit }\begin{array}{lllllllllll:l}{6}&{4}&{4}&{8}&{3}&{2}&{5}&{9}&{7}\\{\hline}
Sort and merge each paired subarray
4 6
of elements
```

Repeat sort/merge until there is only one array

## Merge Sort

Array is divided into its smallest unit
i.e., a single element

Sort and merge each paired subarray | 6 | 4 | 8 | 3 | 2 | 5 | 9 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | of elements

Repeat sort/merge until there is only one array

## Merge Sort

Array is divided into its smallest unit
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Repeat sort/merge until there is only
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46

## Merge Sort



Repeat sort/merge until there is only
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## Merge Sort


i.e., a single element
of elements
Repeat sorke until there is only one array

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## Merge Sort

Array is divided into its smallest unit
i.e., a single element

| 6 | 4 | 8 | 3 | 2 | 5 | 9 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |

Sort and merge each paired subarray

| 4 | 6 | 3 | 8 | 2 | 5 |
| :--- | :--- | :--- | :--- | :--- | :--- | of elements

Repeat sort/merge until there is only
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## Merge Sort

 of elements
Repeat sort/merge until there is only one array

## Merge Sort

| Array is divided into its smallest unit | 6 | 4 | 8 | 3 | 2 | 5 | 9 |  | 7 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| i.e., a single element |  |  |  |  |  |  |  |  |  |
| Sort and merge each paired subarray of elements |  | 6 | 3 | 8 | 2 |  |  |  |  |

of elements
Repeat sort/merge until there is only
one array

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i.e., a single element

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Repeat sort/merge until there is only
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Sort and merge each paired subarray

| 6 | 4 | 8 | 3 | 2 | 5 | 9 | 7 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |


| 4 | 6 | 3 | 8 | 2 | 5 | 7 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | of elements

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Sort and merge each paired subarray of elements
Repeat sort/merge until there is only one array

N.B.: the individual subarrays are already sorted, so we just need to compare the first element in each

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$\square$


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| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 4 | 6 | 3 | 8 | 2 | 5 | 7 | 9 |
| 3 | 4 | 6 | 8 | 2 | 5 | 7 | 9 |
|  | 2 3 4 5 6 7 8 9 |  |  |  |  |  |  |

## Merge Sort Analysis

Called a divide and conquer algorithm
At each level, we look at $n$ elements
Calculating the run time requires also calculating the number of levels
O( $\mathrm{n} \log (\mathrm{n})$ )

both best and worst case

Big O Notation

$\mathrm{O}(\mathrm{n} \log (\mathrm{n}))$ isn't great..
...but it's the best we $\mathrm{O}(\mathrm{n} \log (\mathrm{n}))$ can do for a general sorting algorithm and we've proven that

## Downsides to Merge Sort

Always performs $\mathrm{O}(\mathrm{n} \log (\mathrm{n})$ )
even if the array is already sorted!
Takes up more memory
insertion and selection sort are in-place sorts
i.e., they swap items around in the same array
merge sort requires additional arrays to move from each level

## Quicksort

- Considered one of the classic sorting algorithms
- Similar to merge sort in terms of complexity, run time - another divide-and-conquer algorithm
-Thumbnail sketch:
-repeatedly subdivide elements by comparing to a single element called the pivot - use recursion to sort the subdivisions


## Quicksort

Choose pivot
we'll choose the last element
Subdivide in relation to the pivot
Move pivot
$\begin{array}{llllllll}7 & 4 & 8 & 3 & 2 & 5 & 9 & 6\end{array}$

Sorts subdivisions, repeat until no more divisions can be made

## Quicksort

Choose pivot
we'll choose the last element
Subdivide in relation to the pivot

Move pivot
Sorts subdivisions, repeat until no
more divisions can be made
$\begin{array}{lllllllll}7 & 4 & 8 & 3 & 2 & 5 & 9 & 6\end{array}$

## Quicksort

Choose pivot
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Subdivide in relation to the pivot
Move pivot
Sorts subdivisions, repeat until no
$748 \quad 3 \quad 2 \quad 5 \quad 96$

## Quicksort

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## Choose pivot

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Subdivide in relation to the pivot
Move pivot

Sorts subdivisions, repeat until no | 4 | 3 | 8 | 7 | 5 | 9 | 6 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |

more divisions can be made

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| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |

Sorts subdivisions, repeat until no
more divisions can be made

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Choose pivot
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Subdivide in relation to the pivot
Move pivot
Sorts subdivisions, repeat until no

$$
\begin{array}{|lllll}
4 & 3 & 2 & 5 & 7 \\
7
\end{array} 96
$$

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```
|4
```

more divisions can be made

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Move pivot
Sorts subdivisions, repeat until no

$$
\begin{array}{|lllllllll}
4 & 3 & 2 & 5 & 6 & 7 & 9 & 8 \\
\hline
\end{array}
$$ more divisions can be made

## Quicksort

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$\begin{array}{llllllll}4 & 3 & 2 & 5 & 6 & 7 & 9 & 8\end{array}$
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43020506798
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```
4
```

more divisions can be made

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Move pivot
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Choose pivot
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Move pivot
Sorts subdivisions, repeat until no

| 4 | 3 | 2 | 5 | 6 | 7 | 9 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |

Sorts subdivisions, repeat until no 43 | 2 | 5 | 6 | 7 | 9 | 8 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |

more divisions can be made

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$$
\begin{array}{|l|l|ll|lll|}
\hline 2 & 3 & 4 & 6 & 7 & 9 & 8 \\
\hline
\end{array}
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more divisions can be made

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\hline
\end{array}
$$ more divisions can be made

## Quicksort

## Choose pivot

we'll choose the last element
Subdivide in relation to the pivot
Move pivot
Sorts subdivisions, repeat until no

```
2
```

more divisions can be made

## Quicksort

Choose pivot
we'll choose the last element
Subdivide in relation to the pivot
Move pivot
Sorts subdivisions, repeat until nc

$$
\begin{array}{lllllllll}
2 & 3 & 4 & 5 & 6 & 7 & 9 & 8 \\
\hline
\end{array}
$$

## Quicksort

Choose pivot
we'll choose the last element
Subdivide in relation to the pivot
Move pivot
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What does best case and worst case mean with quicksort?

## Algorithm Analysis

First, consider what is the best and worst case scenarios for sorting an array not the same scenarios for the two different sorts!

Then, fill out the chart below with the run time (i.e., big O ):

|  | merge sort |
| :---: | :---: |
| best case scenario |  |
| worst case scenario |  |
|  | quicksort |
| best case scenario |  |
| worst case scenario |  |

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| :---: | :---: |
| array already sorted <br> (best case scenario) | $\mathrm{O}(\mathrm{n} \log (\mathrm{n})$ ) |
| array sorted backwards <br> (worst case scenario) | $\mathrm{O}(\mathrm{n} \log (\mathrm{n})$ ) |
|  | quicksort |
| pivots are all range midpoints <br> (best case scenario) | $\mathrm{O}(\mathrm{n} \log (\mathrm{n})$ ) |
| pivots are min/max in range <br> (worst case scenario) | $\mathrm{O}\left(\mathrm{n}^{2}\right)$ |

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## Expanding Our Sorting Efforts

What about linked lists? Singly vs doubly linked?
What about objects? How do we define equality/inequality?

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can use any sort that only requires accessing our values in a sequential order
i.e., insertion sort, selection sort, merge sort
quicksort requires random access, and has worse performance
What about objects? How do we define equality/inequality?

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What about linked lists? Singly vs doubly linked?
can use any sort that only requires accessing our values in a sequential order i.e., insertion sort, selection sort, merge sort
quicksort requires random access, and has worse performance
What about objects? How do we define equality/inequality? equals(Object o), compareTo(Object o)

## Merge Sort vs Quicksort, Array vs Linked List

Consider what we know about the strengths and weaknesses of access and insertion in arrays and linked lists. How do those strengths and weaknesses play out in these two sorting algorithms?


